



HeMPS Platform

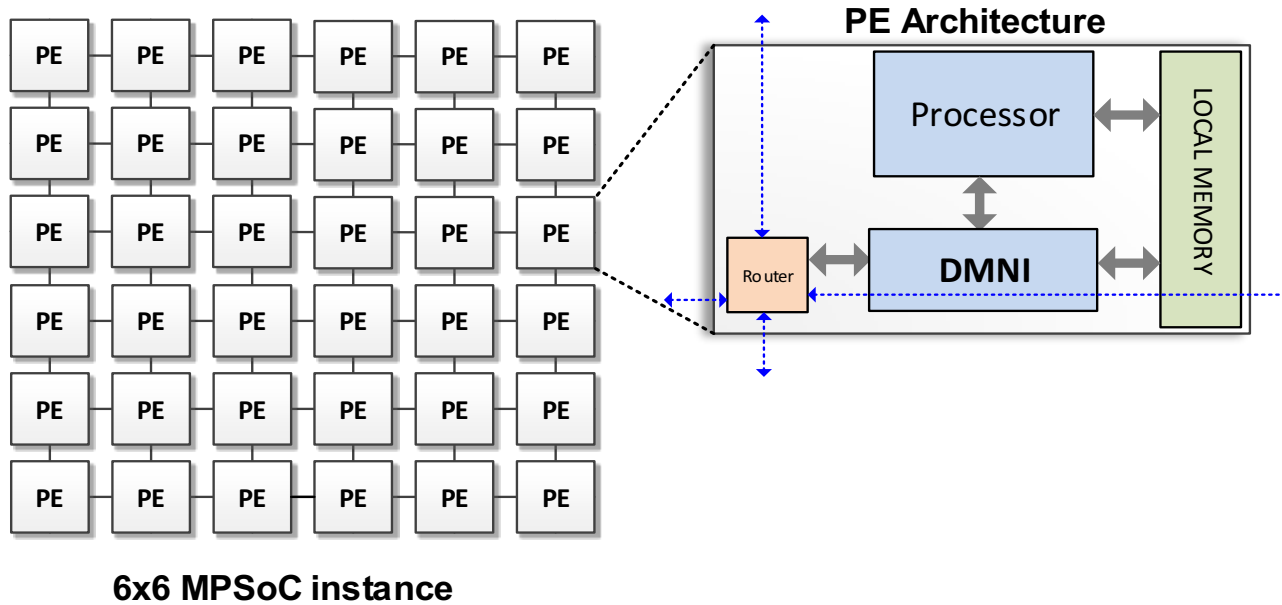
v7.3

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Platform Overview

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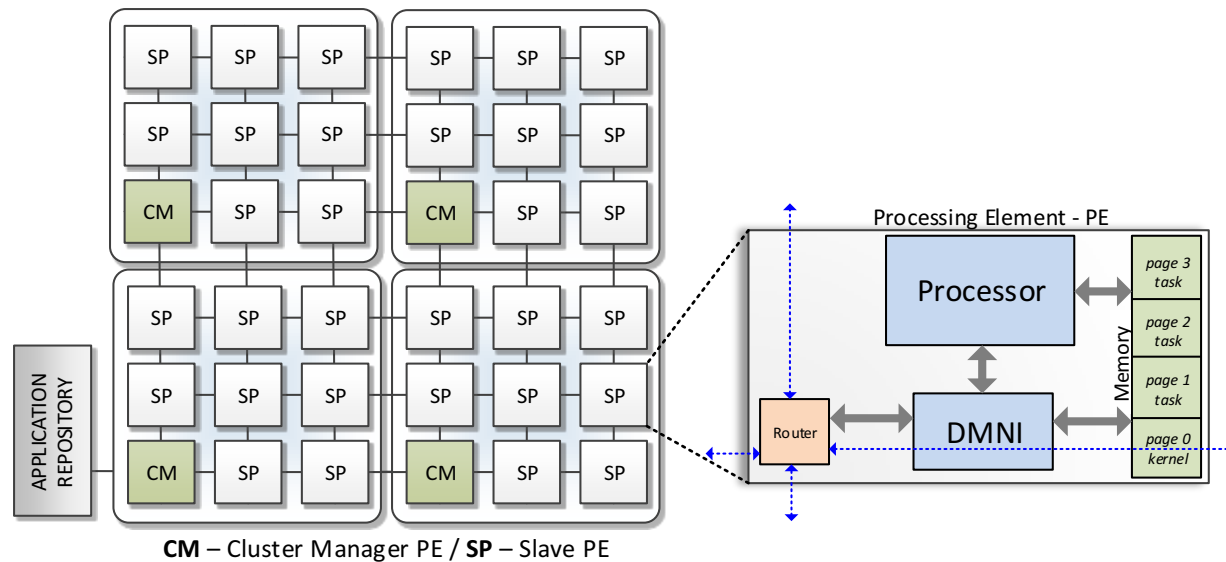
Homogeneous MPSoC

- Each PE has the same architecture

PE is composed of one **processor**, **local memory**, **DMNI**, and **router**

Platform Organization

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Cluster-based organization

- Provides **scalability of management** and **traffic isolation**
- Reclustering is allowed
- Each cluster is managed by a **cluster manager (CM)**
- One CM is responsible for access a external repository containing the application task code

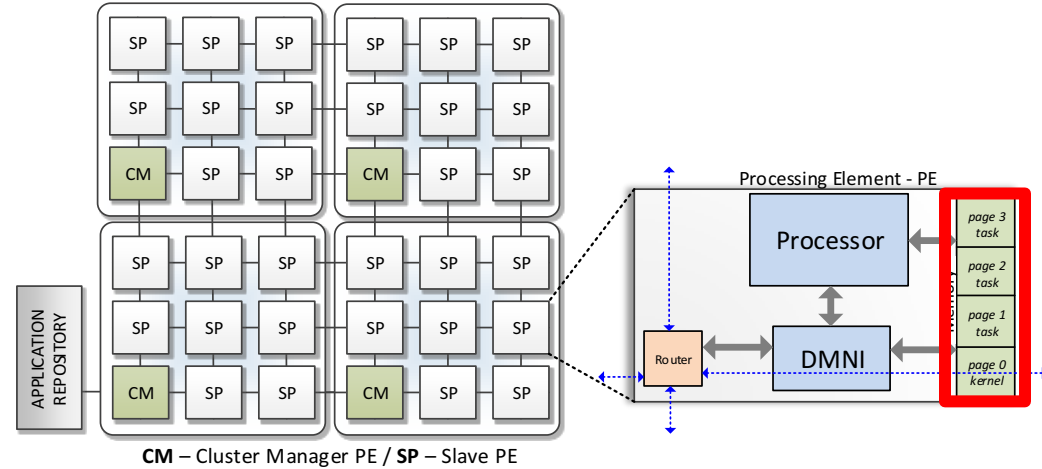
Architectural Features

Local Memory

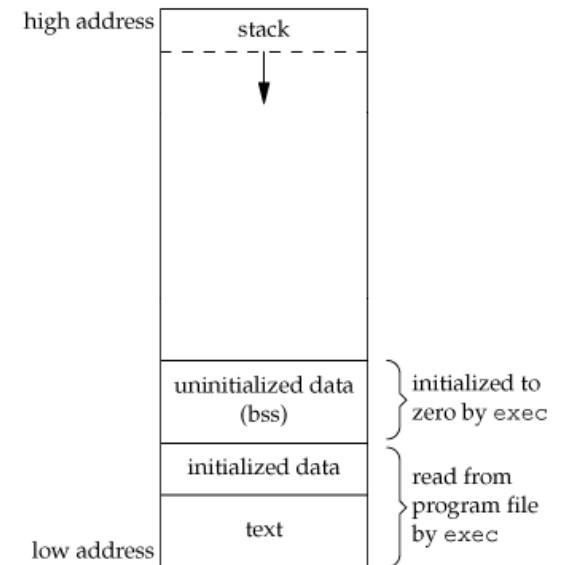
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Scratchpad memory

- RAM
- Dual port
- Size is parameterizable
- Pages are logically managed



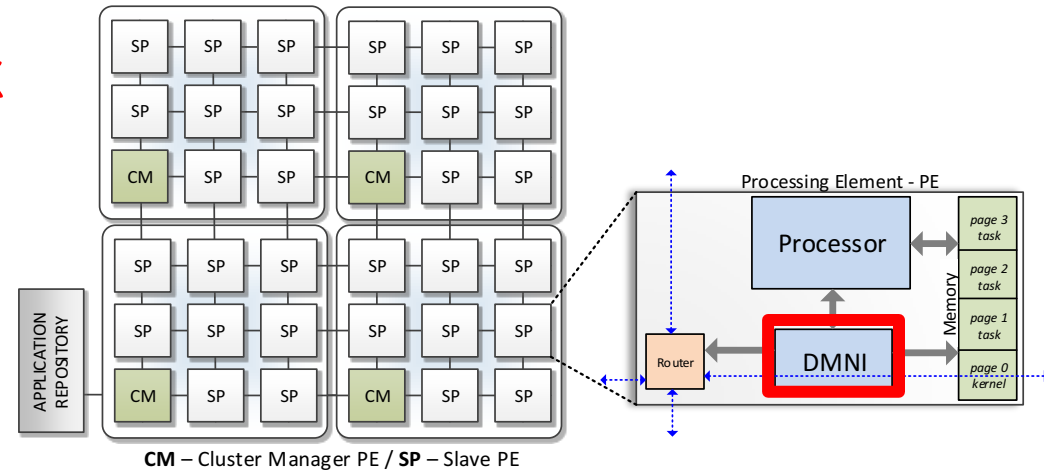
The memory implement a **true dual-port interface** enabling **simultaneous access of processor and DMNI**



DMNI

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Direct Memory Network Interface²



The DMNI implements a **direct interface between the local memory and the NoC**.

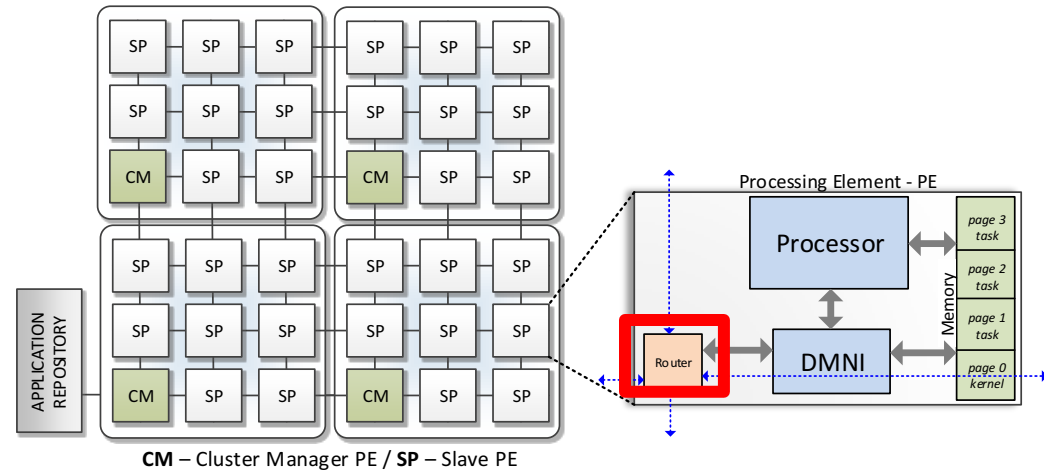
It is an approach specialized to design of NoC-based MPSoC systems

Router

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Hermes NoC³

- XY addressing
- XY and WF routing
- Packet Switching
- Wormhole with credit-based flow control
- Takes 5 clock cycles to arbitrage and routing a packet

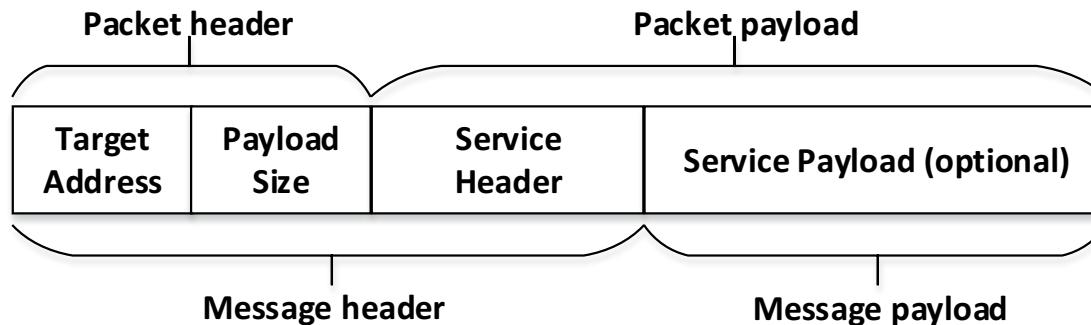


HeMPS v7.3 uses the simplest Hermes NoC implementation. There are several others Hermes derivations

- Asynchronous
- Virtual-channel
- Frequency Scaling
- Circuit-Switching
- Multicast, ...

NoC packet and message structure

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From the NoC point of view, the packet has a **header** and a **payload**

From a task point of view a message contains

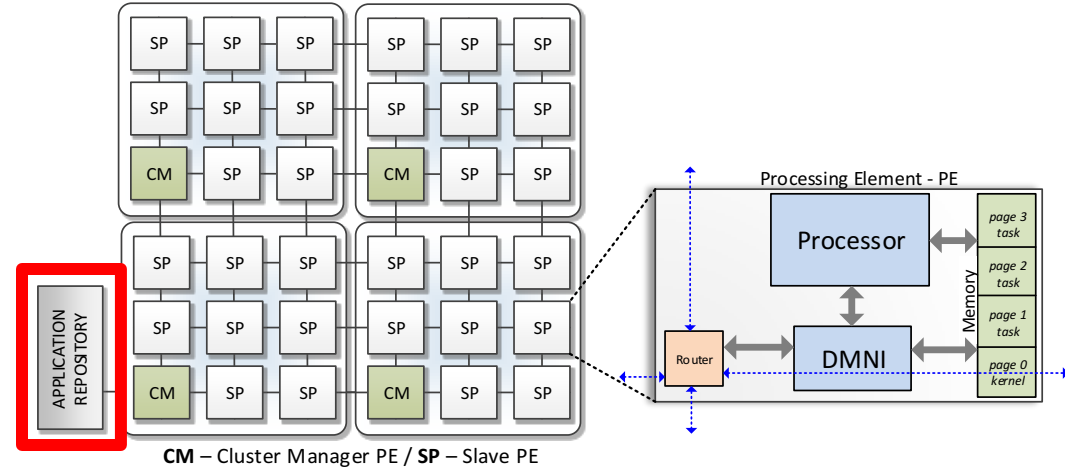
- **Message header**
 - Encapsulates the packet and service header
- **Message payload**
 - Optional field. It may contain for example user data or an object code of a task

Application Repository

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**An external memory
(off-chip)**

Stores the application
description and its task
object code



Logical Features

Logical Features

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Logical Features are implemented by software components

- **μkernels**
 - Slave
 - Manager
- **User's tasks**

Logical Features:

- **System Management**
- **User's Application Execution**

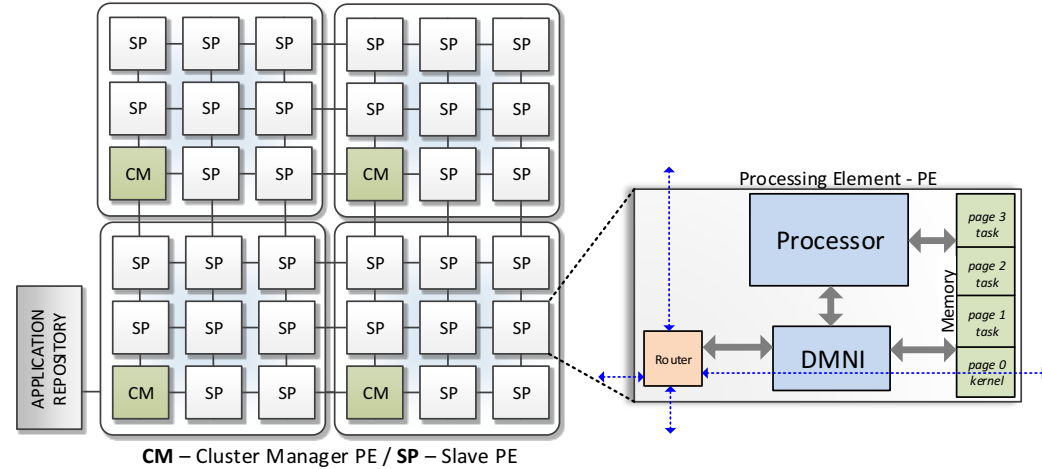
Logical Features

System Management

Cluster-based Management

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Means that the system is logically divided into groups of processors managed by one **Cluster Manager (CM)**



CM is a PE that runs the **manager μ kernel**

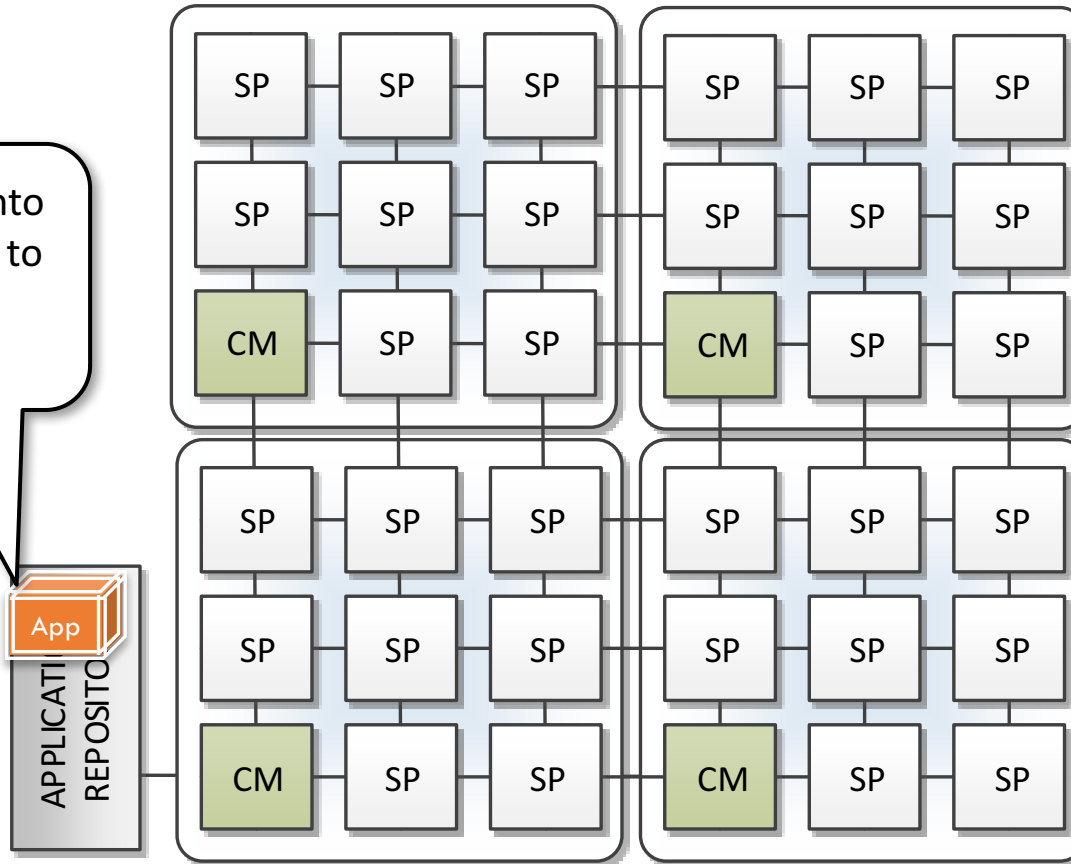
Performs management functions

- Task mapping
- Task migration
- Reclustering

Task Mapping

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App. is stored into AR and request to execute into MPSoC

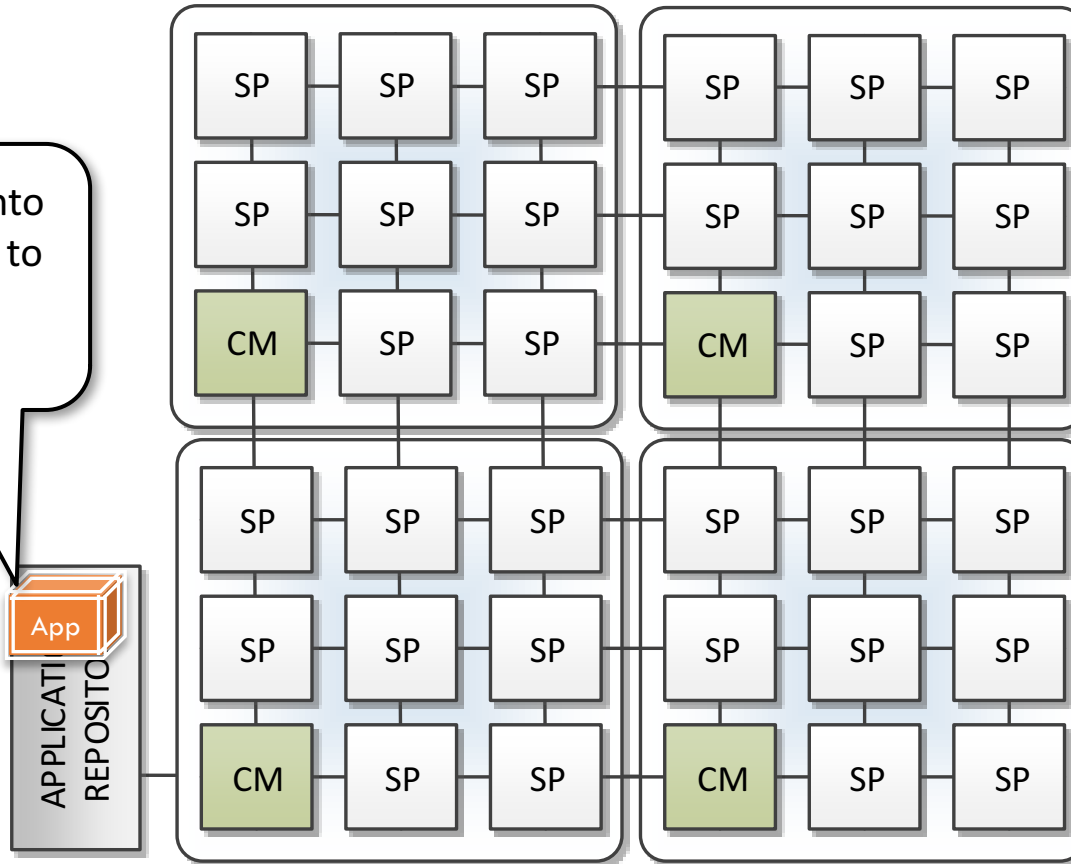


CM – Cluster Manager PE / **SP** – Slave PE

Task Mapping

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App. is stored into AR and request to execute into MPSoC



CM – Cluster Manager PE / **SP** – Slave PE

Task Mapping

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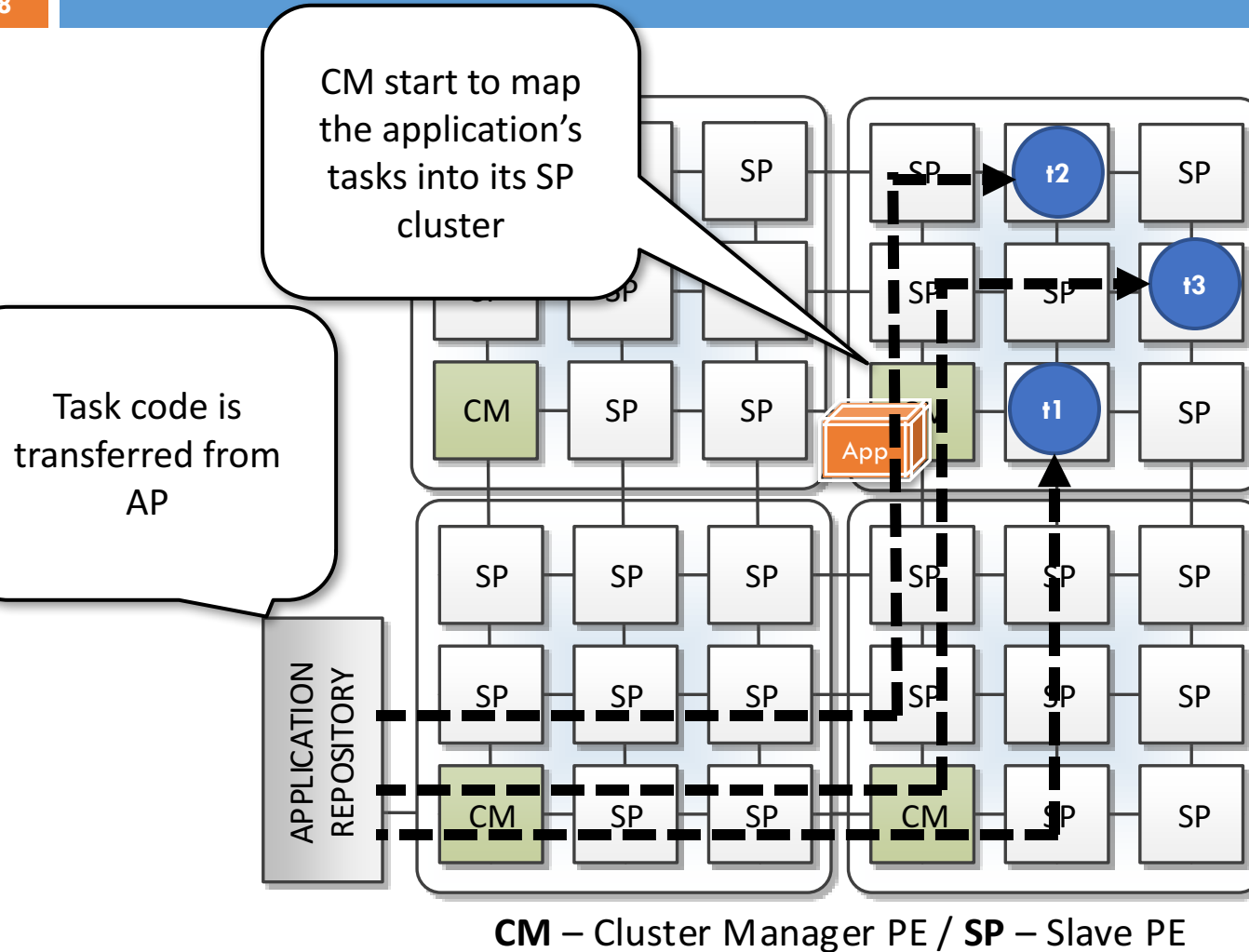
CM handles the request and select an appropriated cluster to receive the App description



CM – Cluster Manager PE / **SP** – Slave PE

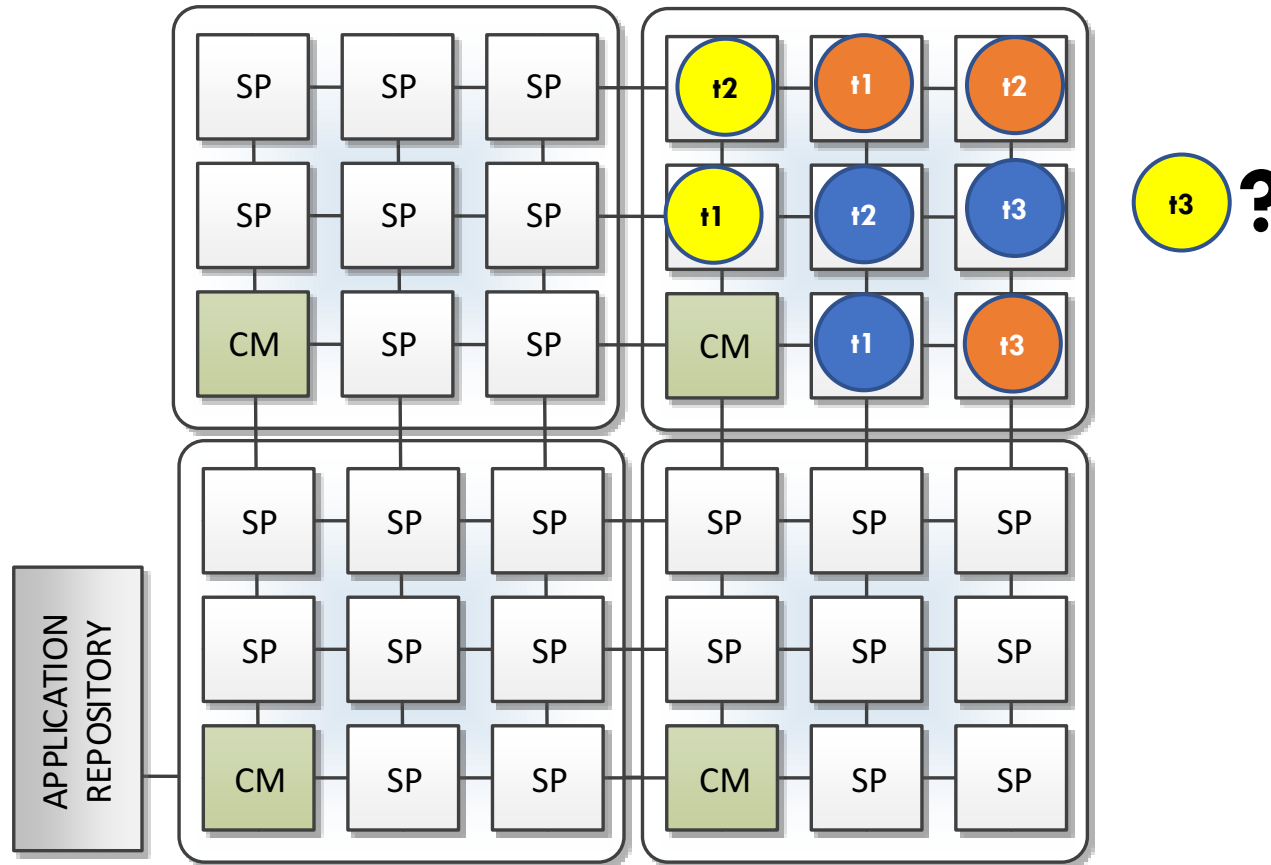
Task Mapping

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Reclustering

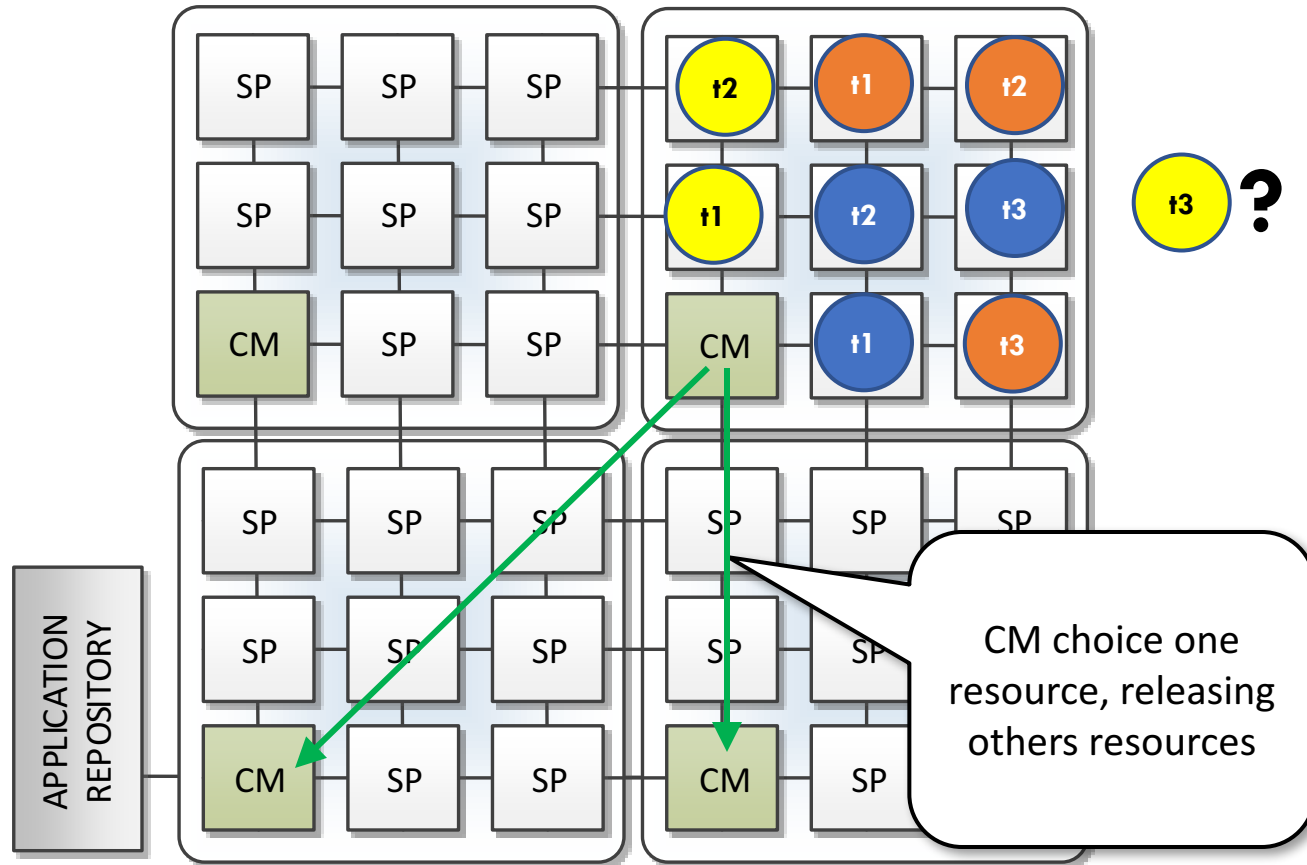
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CM – Cluster Manager PE / **SP** – Slave PE

Reclustering

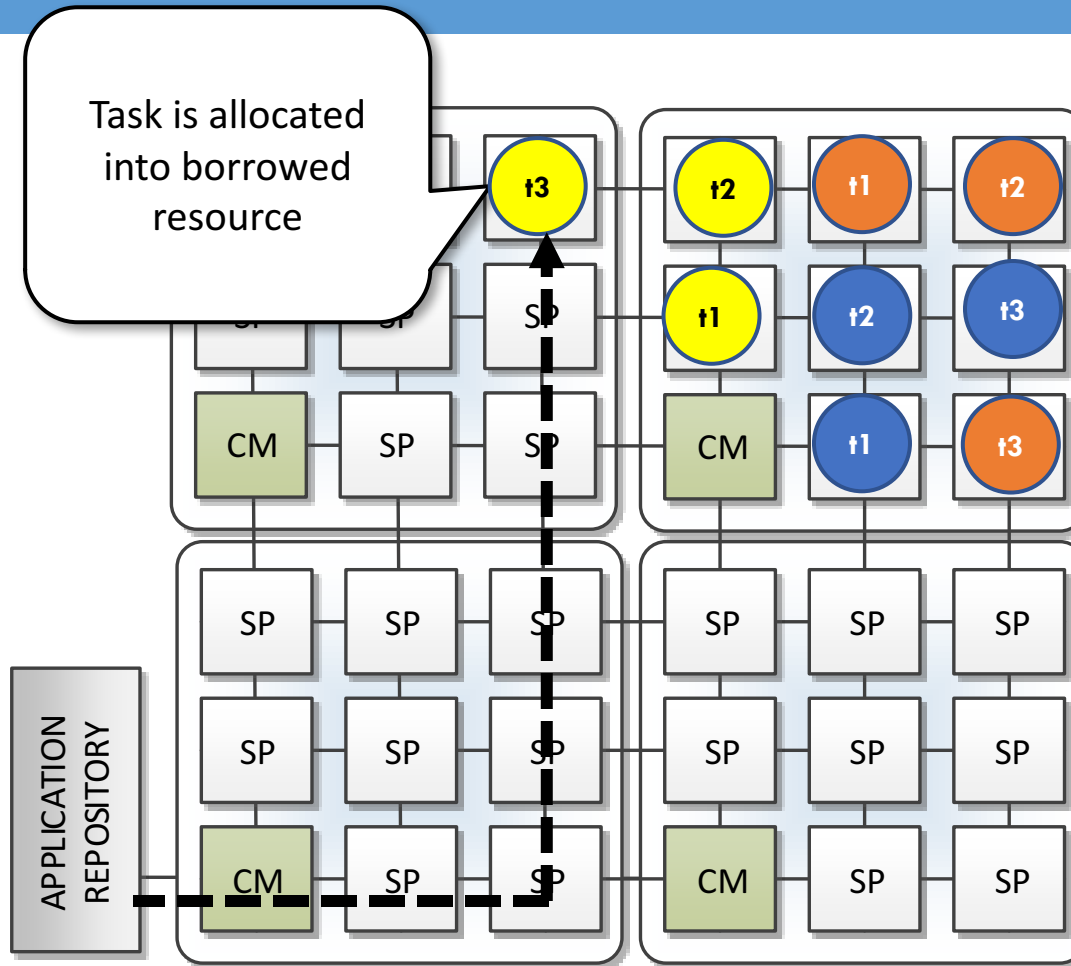
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CM – Cluster Manager PE / **SP** – Slave PE

Reclustering

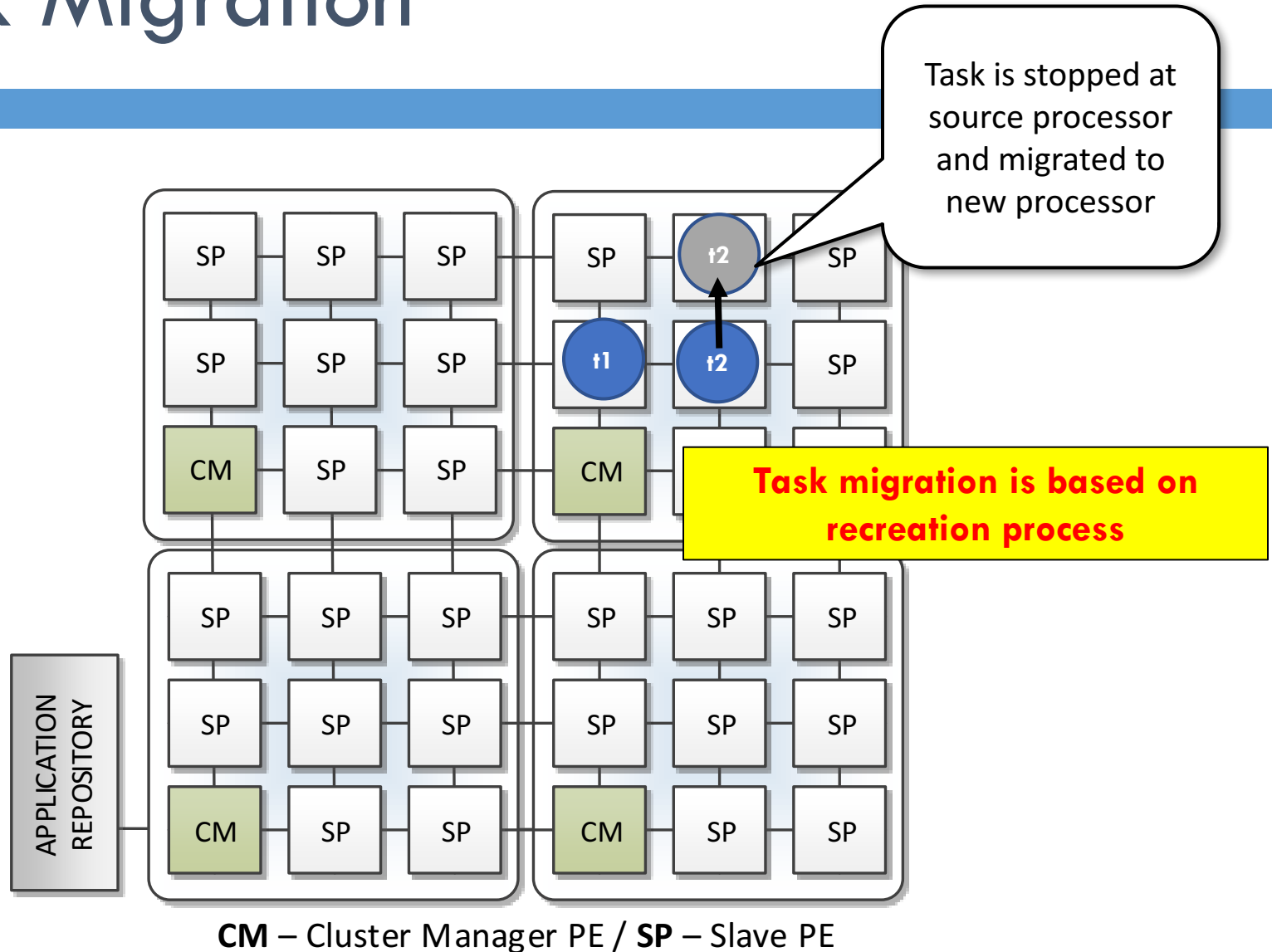
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CM – Cluster Manager PE / **SP** – Slave PE

Task Migration

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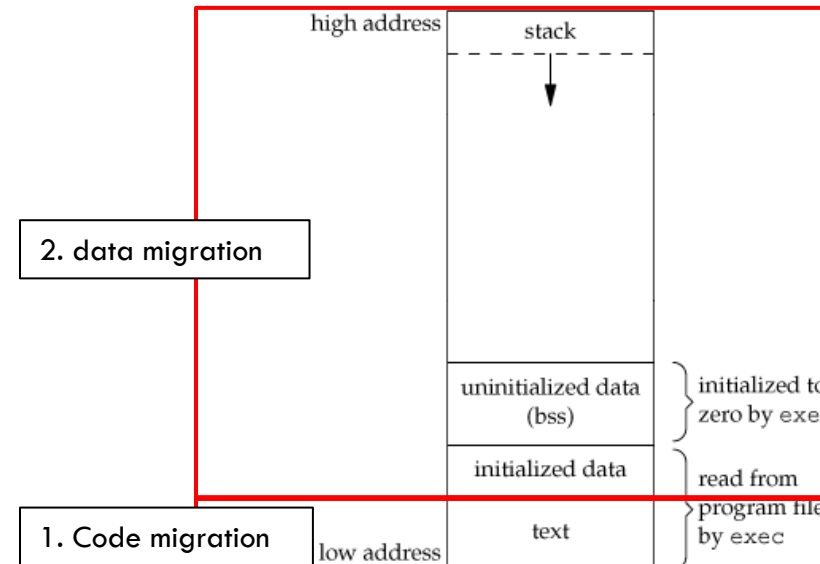
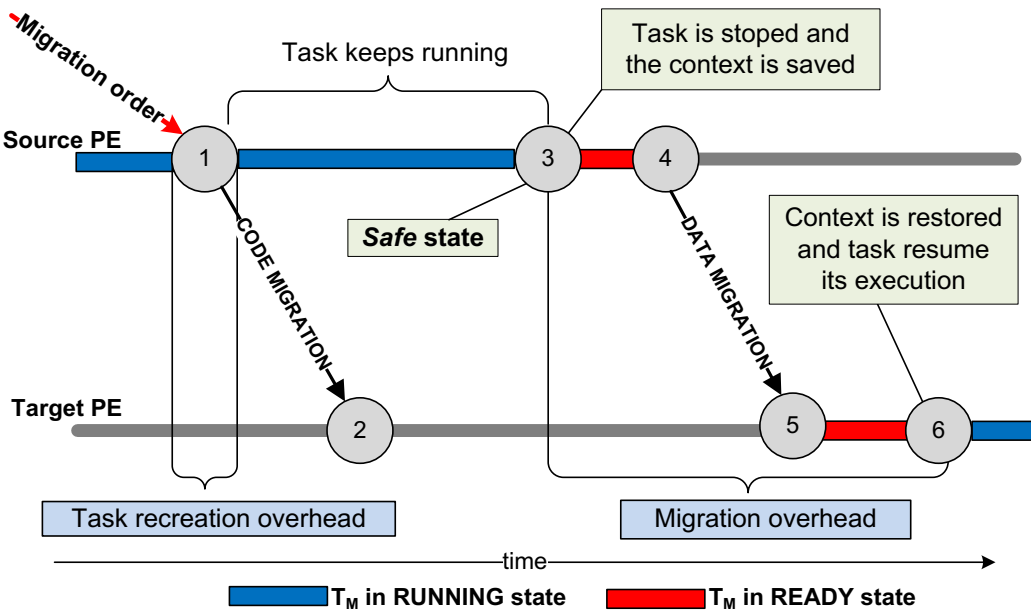


Task Migration

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Migration occurs into steps

- Task **keeps running during its some migration steps**
- Task is only stopped when **safe points** are automatically identified by the migration process (software)
- Safe point are moment which the **task is not waiting for a message from another task**



Logical Features

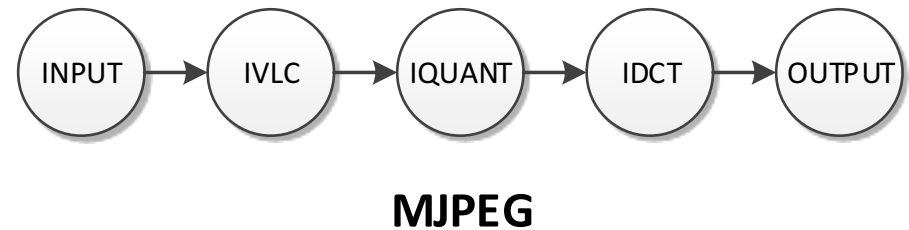
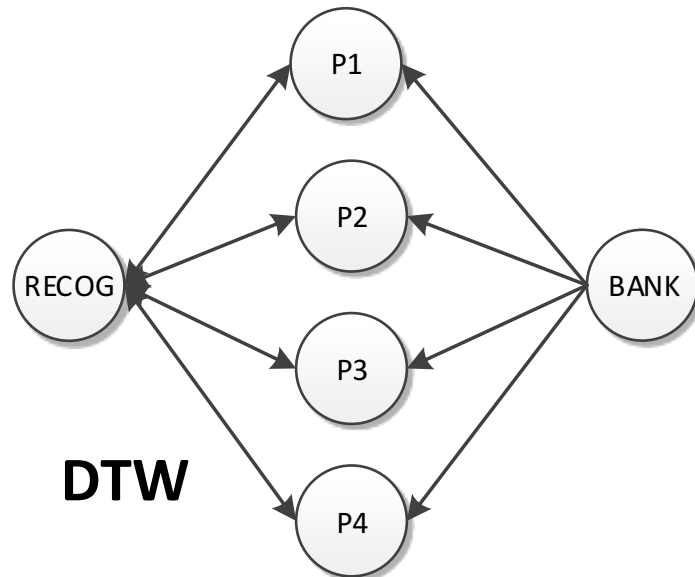
User's Application Execution

Application

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An application is a **set of communicating tasks** (each task is a .c file)

Application are **described as a CTG: Communicating Task Graph**. Example of applications:



Task

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Task is a .c file which perform some computation and communication with other(s) task(s)

Example of a task code

```

int main(){
    int test[SIZE][SIZE];
    int pattern[SIZE][SIZE];
    int result, j;

    Receive(&msg, recognizer);

    Echo("Task P1 INIT\n");

    memcpy(test, msg.msg, sizeof(test));

    for(j=0; j<PATTERN_PER_TASK; j++){

        Echo("Task P1 FOR\n");

        memset(msg.msg,0, sizeof(int)*MSG_SIZE);

        Receive(&msg, bank);

        //Echo("Task P1 received pattern from bank\n");

        memcpy(pattern, msg.msg, sizeof(pattern));

        result = dynamicTimeWarping(test, pattern);

        msg.length = 1;

        msg.msg[0] = result;

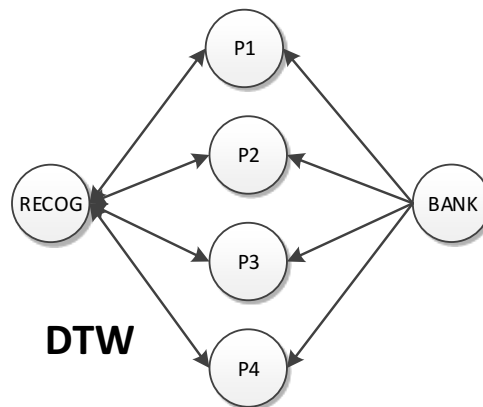
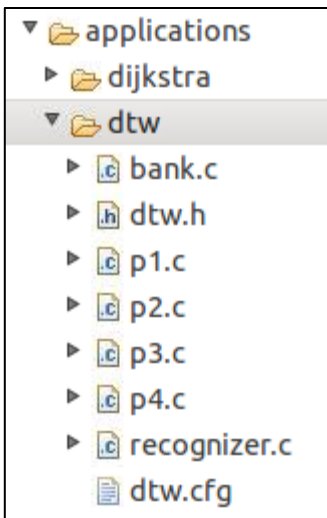
        Send(&msg, recognizer);

    }

    Echo("Task P1 FINISHEDD IN\n");
    Echo(itoa(GetTick()));

    exit();
}
  
```

Example of an application task files



User's Application Execution

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SP are dedicated to execute the user applications

SP is a PE that runs the **slave μ kernel**

Performs support for user task execution

- TCB – Task Control Block
- Inter-task communication
- Scheduling
- Interruption Handling
- API – by System Calls
- Idle



CM – Cluster Manager PE / SP – Slave PE

API

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HeMPS API (MPI-based)

- void **Send**(**Msg *** msg, **unsigned int** target _task_ID)
- void **Receive** (**Msg *** msg, **unsigned int** source _task_ID)
- unsigned int **GetTick**(**void**)
- void **Echo**(**char *** string)
- void **Exit**(**char *** string)

```
typedef struct {
    int length;
    int msg[MSG_SIZE];
} Message;
```

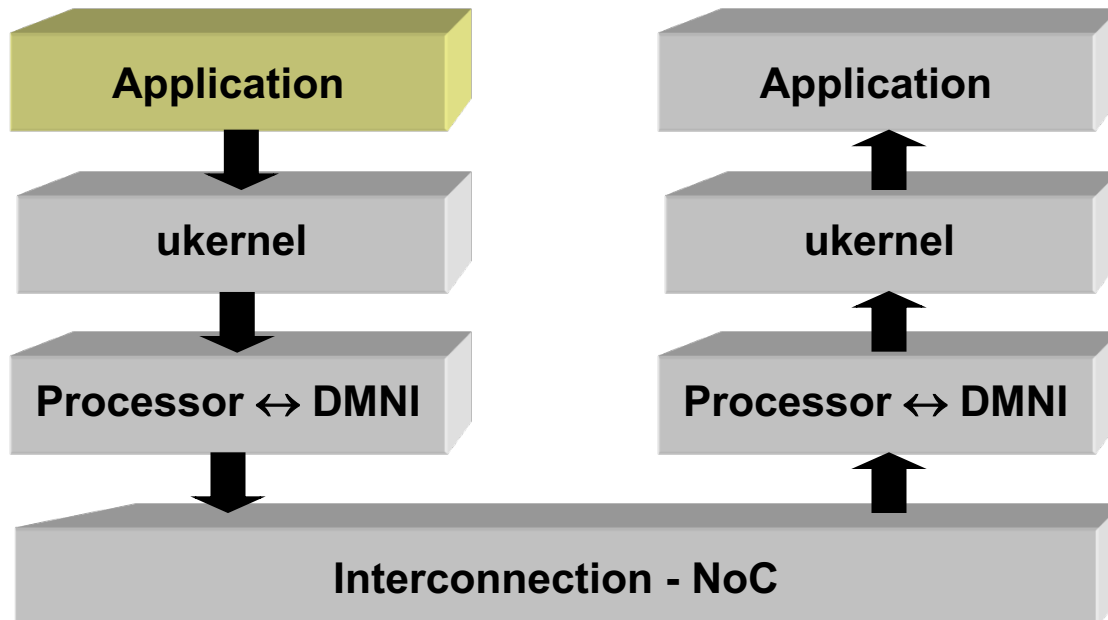
Task communicate using **Send** and **Receive** primitives

Communication Layers

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Application

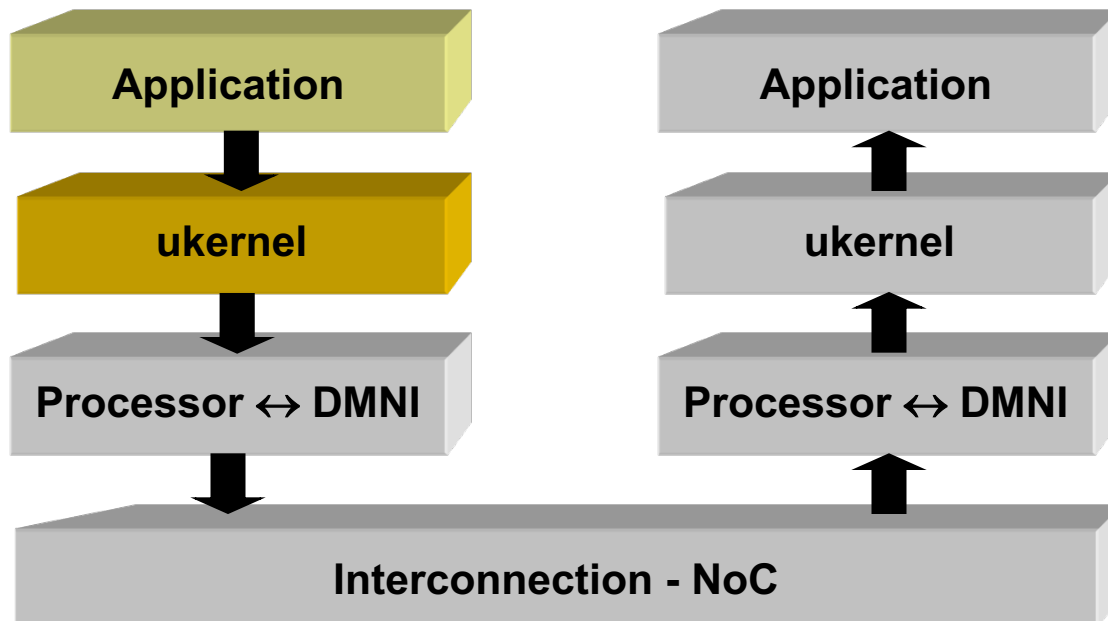
- Send the message by calling the **Send** API primitive



Communication Layers

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ukernel programs the DMNI to **send** memory block in a packet format

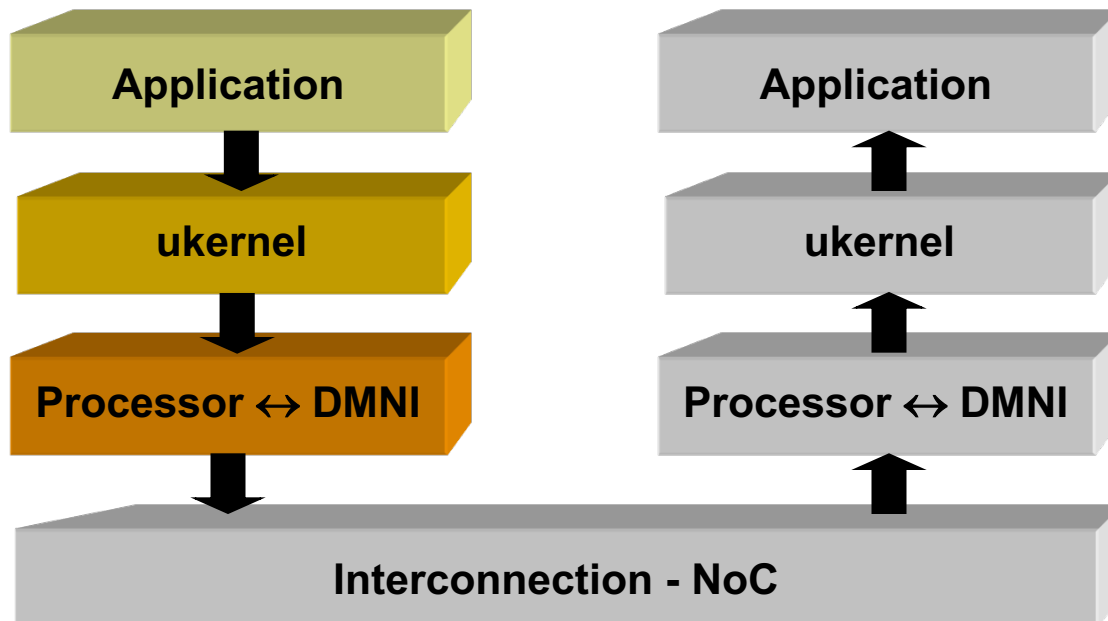


Communication Layers

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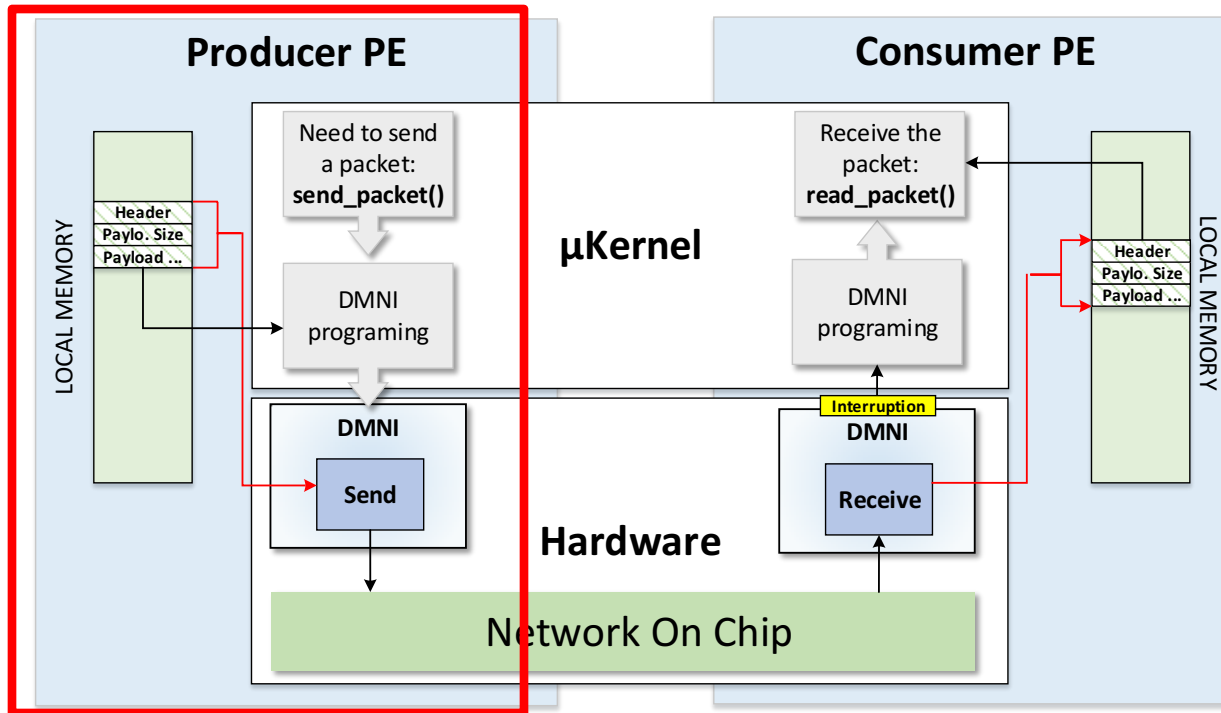
DMNI (send) copies packet from memory and inject into NoC

- Can perform serialization
- Must to implement the NoC flow control



Communication Layers

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send_packet(): function that programs DMNI to copy a memory block to the NoC

- Assumes that the memory block is in the format of a NoC packet

DMNI - Send

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Objective: copy a memory block injecting into the NoC

- The particular feature of this module is the possibility to **transfer two memory blocks with one software programming**

send_packet() API is responsible for to expose the DMNI send feature to the software by configuring MMR

```

1.  void send_packet(mem_size_1, mem_addr_1, mem_size_2,
    mem_addr_2) {
2.      while (MemoryRead(DMNI_SEND_ACTIVE));
3.      MemoryWrite(DMNI_SIZE, mem_size_1);
4.      MemoryWrite(DMNI_ADDRESS, mem_addr_1);
5.      if (mem_size_2 > 0) {
6.          MemoryWrite(DMNI_SIZE_2, mem_size_2);
7.          MemoryWrite(DMNI_ADDRESS_2, mem_addr_2);
8.          MemoryWrite(DMNI_OP, READ);
9.          MemoryWrite(DMNI_START, 1);
10. }
  
```

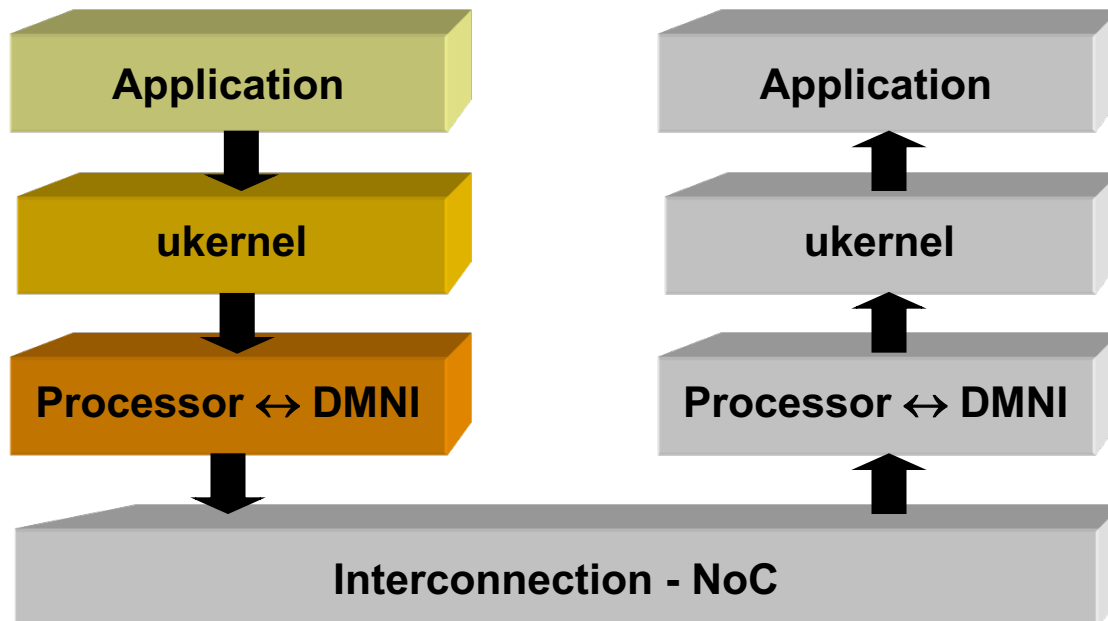
Fig. 5 – *Send_packet()* function, executed in the μ kernel of the processor.

Communication Layers

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NoC – Network on Chip

- Send the packet to the destination PE
- Packet is divided in flits

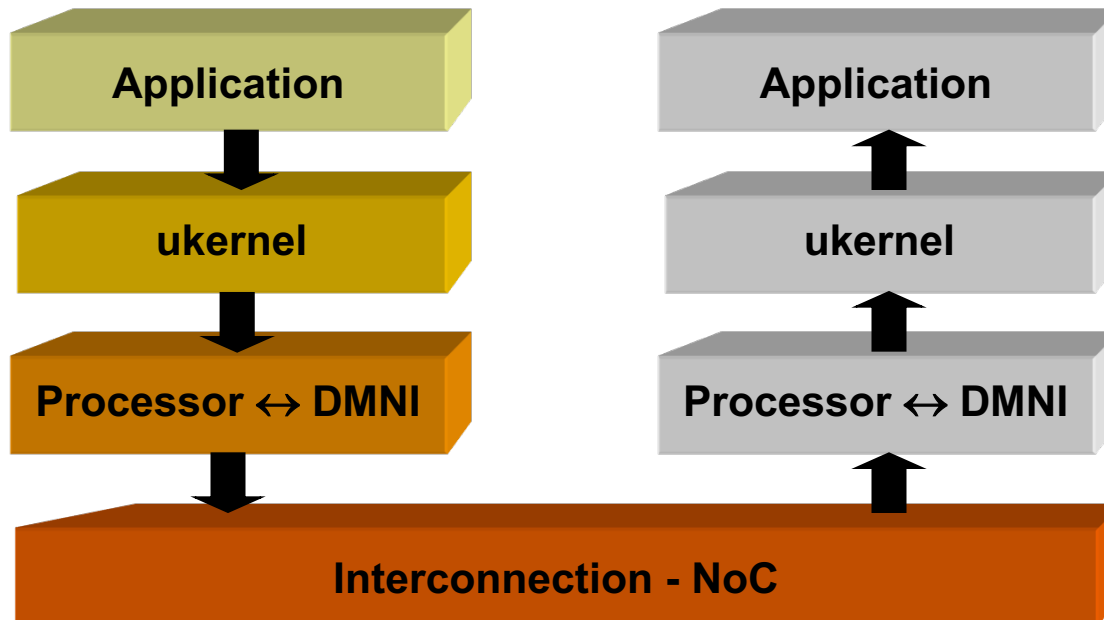
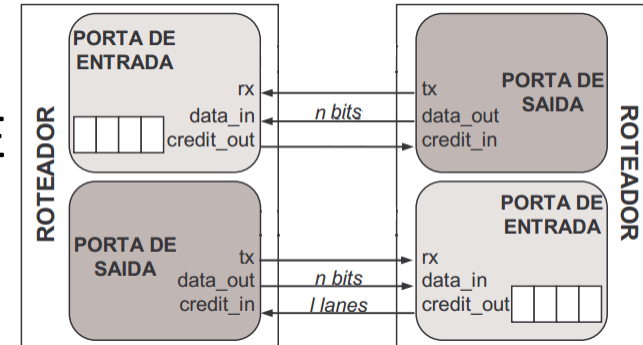


Communication Layers

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NoC – Network on Chip

- Send the packet to the destination PE
- Packet is divided in flits

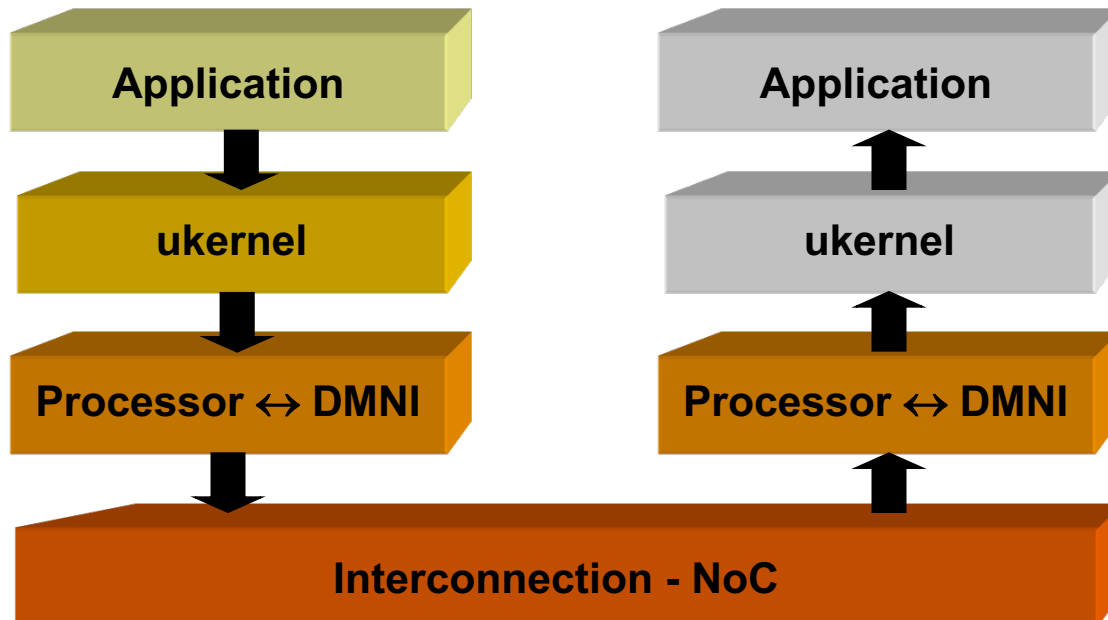


Communication Layers

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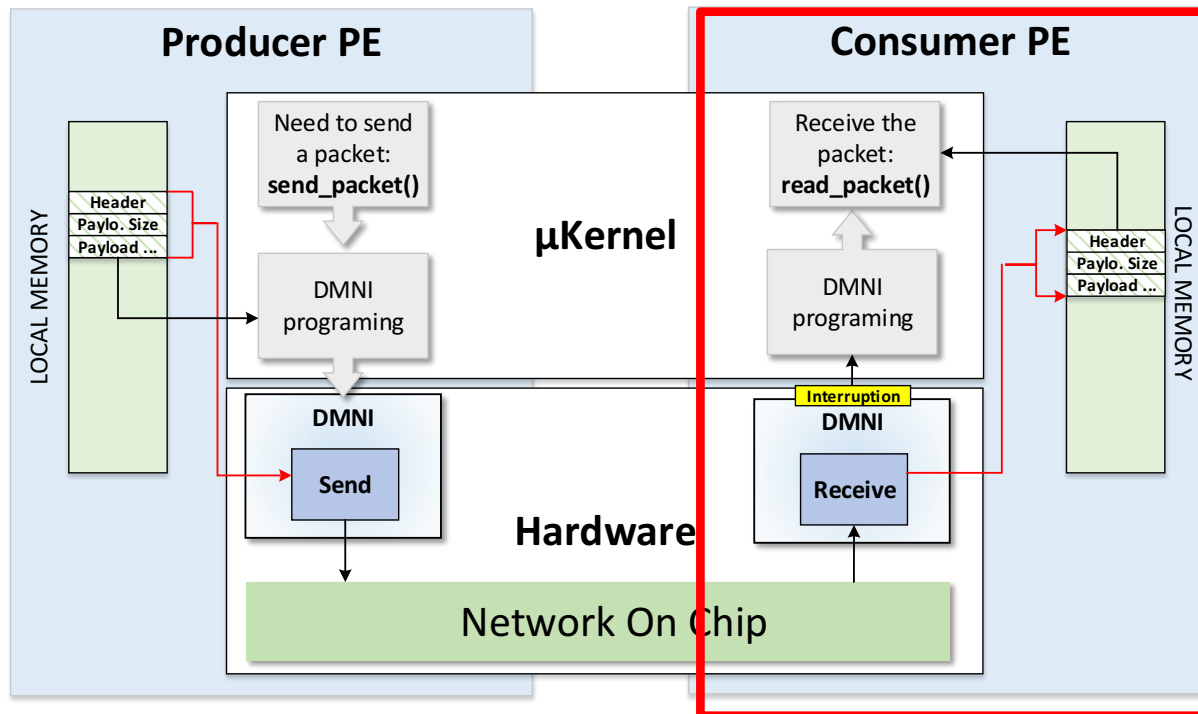
DMNI (receive) copies packet from NoC and transfers into memory

- Can perform deserialization
- Must to implement the NoC flow control



Communication Layers

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read_packet(): function that programs the DMNI to copy a NoC packet to a memory block

- Fired by a interruption

DMNI - Receive

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Objective: receiving a NoC packet coping to a specified memory address

- Also it generates a *software interruption* when detects a incoming packet

receive_packet() API is responsible for expose the DMNI receive feature to the software by configuring MMR

- Called through a *software interruption*, generated when a incoming packet is detected by DMNI

```

1.  void read_packet (init_addr, packet_size)
2.      MemoryWrite(DMNI_SIZE, packet_size);
3.      MemoryWrite(DMNI_ADDRESS, init_addr);
4.      MemoryWrite(DMNI_OP, WRITE);
5.      MemoryWrite(DMNI_START, 1);
6.      while (MemoryRead(DMNI_RECEIVE_ACTIVE));
7.  }
  
```

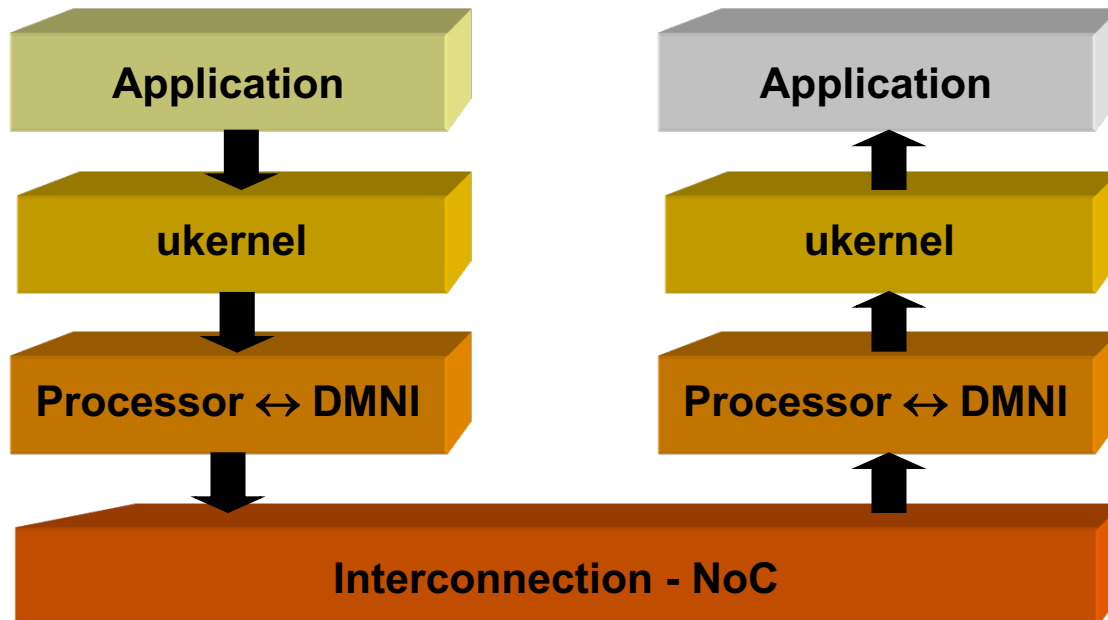
Fig. 9 – *Read_packet()* function, executed in the μ kernel of the processor.

Communication Layers

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ukernel

- Handle packets from DMNI by implementing a interrupt handling mechanism (OS_InterruptServiceRoutine)
- Is responsible to program the DMNI

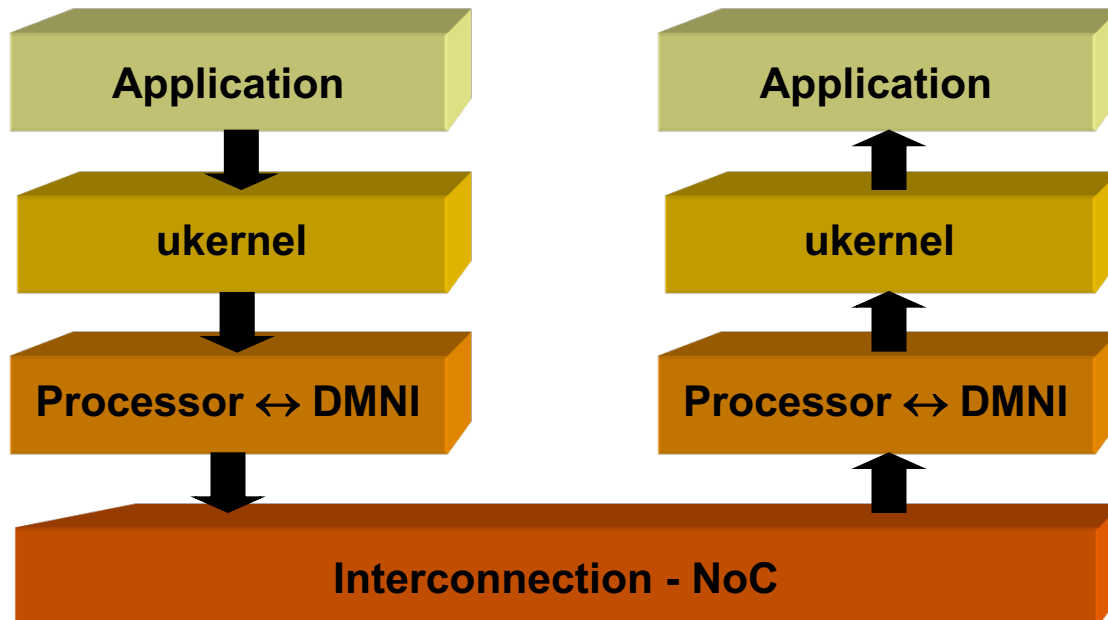


Communication Layers

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Application

- Receive the packet by calling the **Receive** primitive API



Debugging

Debugging

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Debugging can be performed from two perspective

From the system developer viewpoint

- By using the **HeMPS Debugger Tool (HDT)**

From the user viewpoint

- By using **Delorean**
 - Currently integrated into HDT

Debugging Framework

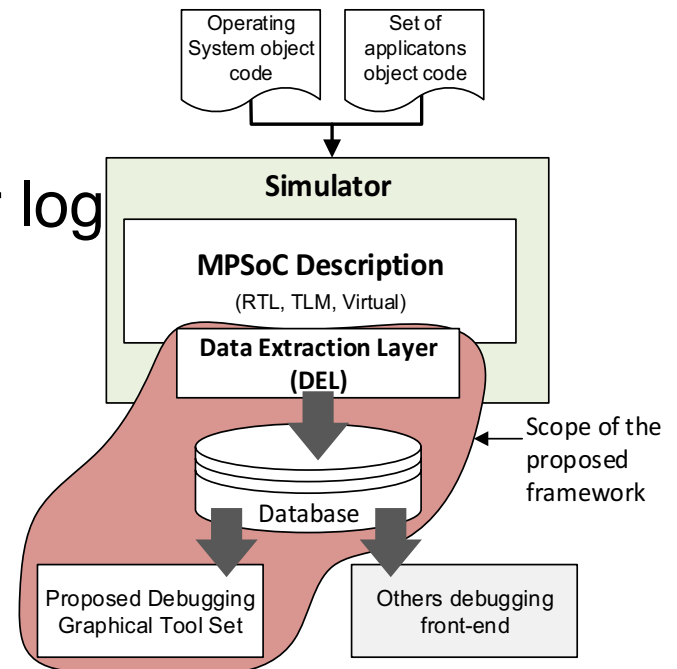
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Data Extraction (back-end)

- Extracts simulated data from platform
- Inserts into a DB or generated log files
- Data extraction following a standard to be generic

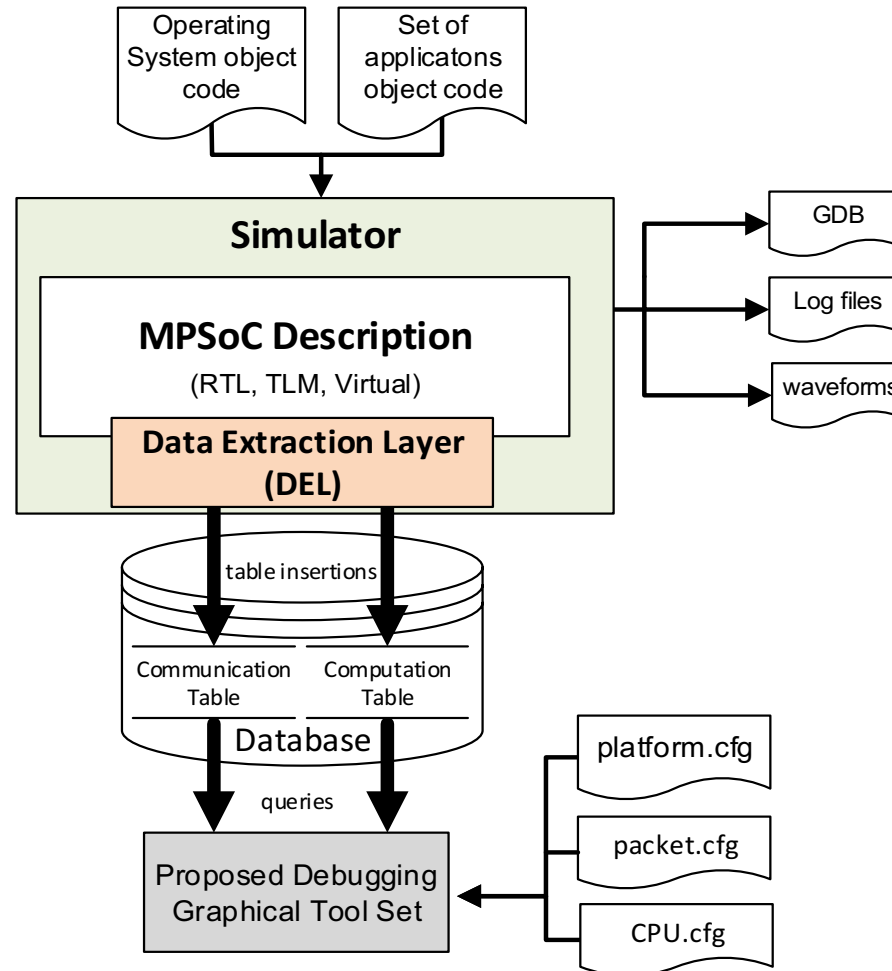
Graphical Debugging (front-end)

- Read extracted data from DB or log files
- Enable easy debugability by the graphical features



Overview

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Main View

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Debug:

- Communication flows
- Routing Algorithms
- Link utilization
- Management Protocols
- Parallel communications



Mapping View

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Debug

- Task mapping algorithm
- PEs occupation
- Task execution status

<input type="checkbox"/> All tasks status <input checked="" type="checkbox"/> Only running <input type="checkbox"/> Only terminated Updating <input type="checkbox"/> Without Task ID			
Slave 0x3 idct 256 RUN iquant 257 RUN	Slave 1x3 start 260 RUN	Slave 2x3 dijkstra_0 512 RUN dijkstra_1 513 RUN	Slave 3x3 dijkstra_4 516 RUN divider 517 RUN
Cluster M 0x2	Slave 1x2 ivlc 258 RUN print 259 RUN	Cluster M 2x2	Slave 3x2 dijkstra_2 514 RUN dijkstra_3 515 RUN
Slave 0x1 bank 768 RUN p1 769 RUN	Slave 1x1 p4 772 RUN recognizer 773 RUN	Slave 2x1 cons 0 RUN prod 1 RUN	Slave 3x1
Global M 0x0	Slave 1x0 p2 770 RUN p3 771 RUN	Cluster M 2x0	Slave 3x0

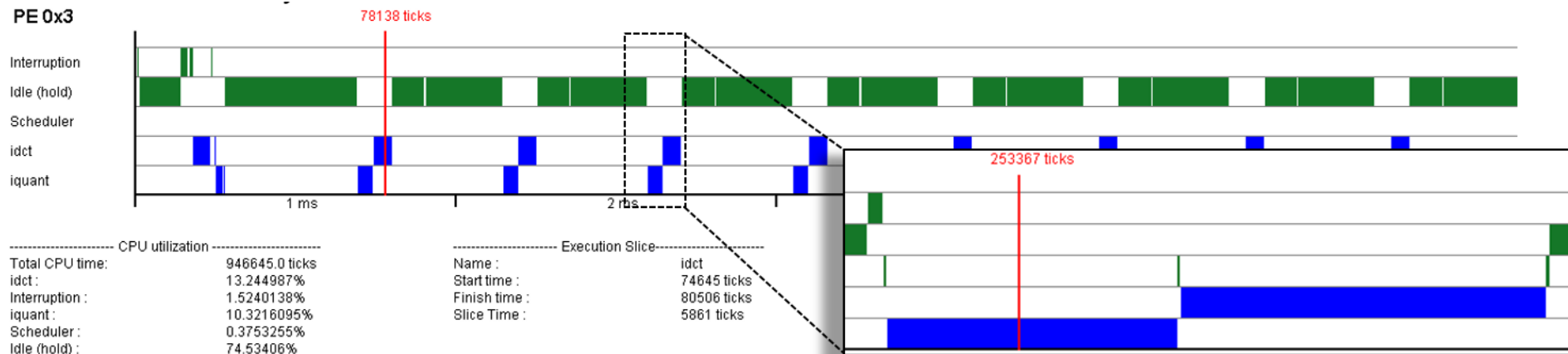
CPU Utilization View

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Enables the designer to verify the CPU use by different software parts over the time

Debug

- Scheduling algorithms
- OS and task bugs
- Other software malfunctions



Delorean

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Used to debug the **application's tasks logs**

It is a **debugging** tool for **the application developer**

